CS 132 - Intro to Computer Science II - Spring 2005 Week #4 Lab Exercise and Homework #3

Week #4 Lab Exercise due: Wednesday, February 9th, END of lab HW #3 due: Wednesday, February 16th, <u>beginning</u> of lab

Week #4 Lab Exercise

INDIVIDUAL-investigation, TEAM-verification exercise:

1. Karen Burgess pointed out a cool site that has animated demonstrations of many searches, including the five discussed in lecture:

http://www.cs.ubc.ca/spider/harrison/Java/sorting-demo.html

On this page, if you click on a picture, you see the corresponding sort in action. If you click on the sort's name, you see this implementation's source code. It is in Java, but the essential sort-logic should still be quite readable for a person who knows C++.

On a sheet of paper with your name on it, answer all of the questions below. Compare them with at least one other classmate's answers. If they differ, discuss why until you both agree on the answer (and change the answer on your paper accordingly if appropriate). Write down the name of all those you conferred with on your paper, then put your name on the Next: list on the board, so that your files can be checked.

- (a) Observe the **bubble sort** demonstration; look at its source code. Can this particular implementation ever stop "early"? If so, how does it tell/determine if it can stop "early"?
- **(b)** Observe the **selection sort** demonstration; look at its source code. Is it selecting and swapping the maximum value from each pass, or the minimum value from each pass?
- (c) Observe the **insertion sort** demonstration; look at its source code. In this implementation, will it always assign the current value-to-be-inserted-next into its proper place in the sorted-portion of the array, even if it is already there (as in the case of an already-sorted array)?
- (d) Note that the **double storage merge sort** here is the rough equivalent to what we discussed in lecture (using a second array to help with the merging); there is also an **in-place merge sort** that performs the merging within the original array (it shifts up the other elements when a merged element needs to go "earlier".) Observe each of these demonstrations, and look at their source code.

Which of these two is faster? Why?

Both of these have the same "base" cases. For what array sizes are there no recursion?

(e) There are four different versions of quick sort here (quick sort, quick sort with bubble sort, enhanced quick sort, and fast quick sort). Observe each of their demonstrations.

The first three use the same method to select their **pivot**. Looking at their source code --- how do these select their pivot (to partition around)?

As described in the reading assignment from the course text (and not in lecture, because of time), partitioning is accomplished by moving from the "left" end until a value is found greater than the pivot, and moving from the "right" end until a value is found less than the pivot, and swapping these two elements. This is then repeated until the two meet in the middle, and that's where the pivot is placed (since everything smaller than the pivot has been swapped to the right of this, and everything larger has been swapped to the left of this). And, the discussion in Ch. 13 of Main/Savitch is *much* better than this quickie description; but, also note that all 4 of these implementations accomplish partitioning in this fashion.

Looking at the first, **quick sort**: for what array sizes is there no recursion? (what array sizes are "base" cases?) Looking at the second, **quick sort with bubble sort**: for what array sizes is bubble sort called instead of quick sort?

To receive credit for this lab exercise, the above must be completed by the end of the lab period.

HOMEWORK #3:

NOTE: you may work TOGETHER on experimentation and analysis for THIS HW (HW #3), but the final WRITE-UP to be TURNED IN must be INDIVIDUAL WORK. (ASK ME if this is not clear!!)

AND NOW, FOR SOMETHING COMPLETELY DIFFERENT... MODIFIED from the assignment by:

David B. Levine, Computer Science Department, St. Bonaventure University

This assignment has the purpose of encouraging you to think more deeply about many of the different sorts that we discussed in Ch. 13 (note that heap sort and radix sort are not included).

Since there are only write-ups involved in this lab, you should turn in the three required components described below ON-PAPER by the beginning of class on the time and date specified above.

On the course web page, via "Homeworks and Handouts", you will find a link named **132hw03.jar**. This file is also available from ~st10/132hw03.jar on cs-server.

One (of several possible) ways to run this is to boot up the NHW 244 lab machine in Linux mode, open up a terminal program, and type the following:

```
cp \simst10/132hw03.jar . # NOTE the period at the end!!! java -classpath 132hw03.jar SortDetective
```

- * You should, of course, let me know if this does not start up the GUI for SortDetective. Note that the window *might* appear underneath other currently-open windows.
- * Also note that the basic approach above might work on any machine that has an appropriate Java JDK installed, changing the directory name appropriately.

SortDetective lets you practice running implementations of **several different** sort implementations, with lists of sizes you choose and various features. It then shows you the **number of comparison** and **number of movements** that that run of that sort took.

The primary objective of this assignment is for you to apply your theoretical knowledge of sorting algorithms to solve a problem of poor user interface design. More specifically, SortDetective is designed to measure comparisons, data movements, and execution time for the several sorting algorithms discussed in class. Unfortunately, the designer of the program did not label the buttons properly. You must apply your understanding of the general properties of the algorithms to determine the proper labeling of the buttons.

The secondary objective of this lab is for you to gain experience writing a concise but complete analysis of a system.

The sorts have dummy names; these are actually implementations of the discussed sorts selection sort, insertion sort, bubble sort, merge sort, and quick sort.

CONSIDER: you know the **basic run-time** complexity of these; you know a bit about how they work. What kinds of tests could you run to make an **educated guess** about which sort is which?

As you know from class, if you double the size of the data set that you give to a quadratic algorithm, it will do four times the work; by contrast, an O(NlogN) algorithm will do a bit more than twice as much; and, a linear algorithm will do only twice as much work. As you also know, the characteristics of the input data set can affect the expected performance of many of our sorting algorithms. Before you begin the planning, you can see how it would be helpful for you to review the **expected** performance of the algorithms on various data sets.

FIRST ITEM TO TURN IN:

First, make a plan --- what kinds of tests could you run to GUESS which sort is which? Write this plan/strategy out, and turn it in as part of this assignment. (You may discuss potential plans/strategies with one another, but each should

write out their plan individually.) Note the this plan should be word-processed rather than hand-written.

What should this plan/strategy look like? Think of the procedure for a science experiment --- if you gave this to two different people, both could go through these steps and be able to compare their results. That is, it should be specific! What sizes of lists will you run the different sorts for? What characteristics will these lists have? What data will you then record for each test? And, you should also include your hypotheses for what you hope to discover for each set of tests.

<u>For example</u>, if you were making a plan to try to distinguish sequential search from binary search, you might have a plan such as:

* Run each search on a list of size 32 for an element not in the list. Record the number of comparisons.

Repeat for 5 different non-list elements: (one smaller than any in the list, one larger than any in the list,

one near the center of the "range" of the list, one near the beginning of the "range" of the list, one near the end of the "range" of the list)

Hypotheses:

- * the binary search should take about the same number of steps for each non-list element.
- * the sequential search may vary [depending on what "shortcuts" are built in --- you'd probably know what these were at the time you are making the plan. You will know for the sorts, as you will see below.]
- * Run each search on a list of size 64 for an element not in the list. Record the number of comparisons.

Repeat for 5 different non-list elements: (one smaller than any in the list,

one larger than any in the list,

one near the center of the "range" of the list, one near the beginning of the "range" of the list, one near the end of the "range" of the list)

Hypotheses:

- * the binary searches here will also take about the same number of comparisons for each search, and should only be a few more steps than the size-32 searches took.
- * the sequential searches here may vary, but at least some should take roughly twice as many steps as the corresponding size-32 searches.

SECOND ITEM TO TURN IN:

Then, **execute** your plan, running the tests and **recording** the results. (You may hand-record these results, rather than word-process or type them, if you wish.) (Yes, we are practicing a little scientific experimentation here, too.) You may discuss and brainstorm together during this testing, also.

Note that you will likely have a number of tables/charts as the result of executing your plan and recording the results (the results will be recorded in tables or charts, which, again, may be either hand-written or typed).

What if you find you need additional tests? Run them, but record for each what you did, and why you did them, along with their results.

THIRD ITEM TO TURN IN:

Finally, you will take your notes from those experiments, and write up your guesses as to which sort is which. You will first summarize your guesses, and then justify them; the format should be like that given on the example on the next page. (Although you may experiment, speculate, and brainstorm with other class members, THIS write-up is INDIVIDUAL work) Like the original plan, this final report should be word-processed rather than hand-written.

Searching Lab Report A. Beauregard Clump 2-14-05

General Results

Button Name Searching Algorithm

Sigma Binary Search
Tau Sequential Search

Rationale

I chose to search data sets of size 32 and 64, looking for 5 elements not in the list each time: one smaller than any in the list, one larger than any in the list, one that is about in the middle of the "range" of elements in the list, one that is near the beginning of the "range" of elements in the list, and one that is near the end of the "range" of elements in the list. Button Tau took very different numbers of steps for the 5 different searches in the 2 lists, but the largest numbers of its comparisons were doubled for the size-64 list as compared to the size-32 list. Button Sigma was very consistent in its number of comparisons for each item in the size-32 list, and this number increased by only 2 for the size-64 searches.

Since a binary search is logarithmic, and the number of comparisons for Tau only increased by a small amount when the list was doubles, I conclude that Button SIGMA is BINARY SEARCH.

Since a sequential search is linear in the worst and average cases, and the highest-comparison searches did double in number of comparisons when the list size was doubles, I conclude that Button TAU is SEQUENTIAL SEARCH.

GRADING NOTES:

- * There is no coding in this assignment. Thus, you should expect that a significant portion of the grade for this assignment will be determined by the quality of the writing of the report. This includes the completeness of the report, the clarity of the writing, and general presentation. You could get a poor grade for poor writing, even if you "get" all of the right matches.
- * Some of the sorts may be difficult to distinguish. A carefully outlined experiment may compensate for an error in these cases if the writing makes it clear that your conclusions/guesses are substantiated by the data.
- * Remember that your report (the 3rd item you are turning in) needn't detail every experiment you ran. Rather, it should give sufficient information to justify your conclusions. It is possible to write a very short report that is completely correct if your experiments are well-chosen. After you determine the matching, you might consider whether there was a shorter way to arrive at your conclusion!
- * Carefully note the requirements given; not turning in some required pieces, or structuring them differently, etc., could have a negative impact on your grade.

REFERENCES:

Below, you will find the Java implementations of the sorts used within the SortDetective program. Again, these should be readable by someone with a knowledge of C++. Why might you want to look at the implementations? To answer such questions as, does this bubble sort stop "early" if it can? How does this quick sort choose its pivot? Does this merge sort use a temporary "helper" array for the merge step? Does this insertion sort always recopy the current element into its insertion location in the sorted part of the array, even if it is already there? etc.

CS 132 - Homework #3 Spring 2005

}

}

comparisons++;
} while (temp != list[0]);

/* Here are the implementations of the sorts as used in SortDetective. *//st (note that they have had code added to count comparisons and data */ /* movements) /* Code and Comments by David B. Levine, Computer Science Department, *//* St. Bonaventure University */ /* BubbleSort */ // Code basically taken from Sedgewick, "Algorithms" protected void bubbleSort(int list[]) int temp; do { movements++; temp = list[0];for (int j=1; j<list.length; j++)</pre> comparisons++; if (list[j-1]>list[j]) { movements += 3; temp = list[j]; list[j] = list[j-1];list[j-1] = temp;}

```
/* Selection Sort */
// Code from memory
protected void selectionSort(int list[])
         for (int j=list.length-1; j>0; j--)
                int maxpos = 0;
                for (int k=1; k <= j; k++)
                      comparisons++;
                      if (list[k]>list[maxpos])
                            maxpos = k;
                }
                if (j != maxpos) // Only move if we must
                      movements += 3;
                      int temp = list[j];
                      list[j] = list[maxpos];
                      list[maxpos] = temp;
                }
}
/* Insertion Sort */
// Code basically taken from Sedgewick, "Algorithms"
protected void insertionSort(int list[])
         for (int j=1; j<list.length; j++)</pre>
                movements++;
                int temp = list[j];
                int k = j;
                while ( k > 0 \&\& list[k-1]>temp )
                      comparisons++;
                      movements++;
                      list[k] = list[k-1];
                      k--;
                if (k > 0)
                {
                      comparisons++;
                movements++;
                list[k] = temp;
         }
}
```

/* MergeSort */

```
// Code basically taken from Horowitz and Sahni, "Fundamentals of Computer
// Algorithms"
protected void mergeSort(int list[])
{
            msort(list, 0, list.length-1);
private void msort(int list[], int low, int high)
            if (low<high)
            {
                    int mid = (low+high)/2;
                    msort(list, low, mid);
                    msort(list, mid+1, high);
merge(list, low, mid, high);
            }
}
private void merge(int list[], int low, int mid, int high)
{
            int h = low;
            int i = 0;
            int j = mid+1;
            int otherList[] = new int[high-low+1];
            while ((h <= mid) && (j <= high))
            {
                     comparisons++;
                    if (list[h] <= list[j])</pre>
                     {
                            movements++;
                             otherList[i] = list[h];
                             h++;
                     }
                     else
                     {
                             movements++;
                             otherList[i] = list[j];
                             j++;
                     i++;
            if (h>mid)
                     for (int k=j; k \le high; k++)
                            movements++;
                            otherList[i] = list[k];
                             i++;
                     }
                     else
                     {
                             for (int k=h; k \le mid; k++)
                                     movements++;
                                     otherList[i] = list[k];
                                     i++;
                             }
                     }
                     for (int m=0; m<otherList.length; m++)
                     {
                            movements++;
                            list[low+m] = otherList[m];
                     }
         }
}
```

}

/* QuickSort */ // Code basically taken from Cormen, Leiserson, and Rivest, "Introduction to // Algorithms" protected void quickSort(int list[]) qsort(list, 0, list.length-1); private void qsort(int list[], int low, int high) if (low<high) int pivot = partition(list, low, high); qsort(list, low, pivot); qsort(list, pivot+1, high); } private int partition(int list[], int low, int high) movements++; int temp = list[low]; int i = low-1; int j = high+1;while(true) do { j--; comparisons++; } while (list[j] > temp); do { i++; comparisons++; } while (list[i] < temp);</pre> if (i < j){ movements += 3; int swapTemp = list[i]; list[i] = list[j]; list[j] = swapTemp; else { return j;